



Fall 2015 COMP 3511

Operating Systems



Lab 05

Outline

- Review Questions
- Nachos Threads and Example
- Nachos Thread Scheduling
- Nachos Thread Switching

Review Questions

- PPT version

- http://course.cs.ust.hk/comp3511/lab/lab05/lab05_review.ppt

- PDF version

- http://course.cs.ust.hk/comp3511/lab/lab05/lab05_review.pdf

Nachos Threads

- In Nachos (and many systems) a ***process*** consists of:
 - An ***address space***, which is further broken down into:
 - 1) Executable code
 - 2) Stack space for local variables
 - 3) Heap space for global variables and dynamically allocated memory
 - A single thread of control, e.g., the CPU executes instructions sequentially within the process
 - Other objects, such as open file descriptors

Nachos Threads

- It is sometimes useful to allow multiple threads of control to execute concurrently within a single process.
- These individual threads of control are called *threads*.
- **One big difference** between threads and processes is that global variables are shared among all threads.

Nachos Threads

- Nachos provides *threads*
 - Nachos threads execute and share the same code (the Nachos source code)
 - and share the same global variables
- The Nachos *scheduler* maintains a data structure called a *ready list*, which keeps track of the threads that are ready to execute.

Nachos Threads

- Threads on the ready list are ready to execute and can be selected for executing by the **scheduler** at any time.
- Each thread has an associated **state** describing what the thread is currently doing.
- Nachos' threads are in one of four states: **READY**, **RUNNING**, **BLOCKED**, **JUST_CREATED**

Nachos Threads

■ **READY:**

- The thread is eligible to use the CPU (e.g., it's on the **ready list**), but another thread happens to be running.
- When the scheduler selects a thread for execution, it removes it from the ready list and changes its state from **READY** to **RUNNING**.
- Only threads in the **READY** state should be found on the **ready list**.

Nachos Threads

■ **RUNNING:**

- The thread is currently running.
- **Only one thread** can be in the **RUNNING** state at a time.
- In Nachos, the global variable ***currentThread*** always points to the currently running thread.

Nachos Threads

■ **BLOCKED:**

- The thread is blocked waiting for some external **event**; it cannot execute until that event takes place.
- Specifically, the thread has put itself to sleep via *Thread::Sleep()*.
- It may be waiting on a condition variable, semaphore, etc.

Nachos Threads

■ **JUST_CREATED:**

- The thread exists, but has no stack yet.
- This state is a temporary state used during thread creation.
- The *Thread* constructor creates a thread, whereas *Thread::Fork()* actually turns the thread into one that the CPU can execute (e.g., by placing it on the **ready list**).

Nachos Threads

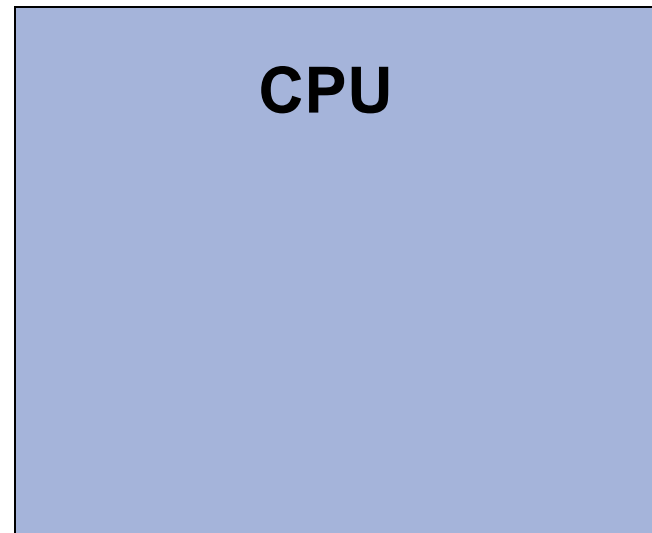
- Nachos does not maintain an explicit process table.
- Instead, information associated with thread is maintained as (usually) private data of a ***Thread*** object instance.
- To get at a specific thread's information, a pointer to the thread instance is needed.

Nachos Threads

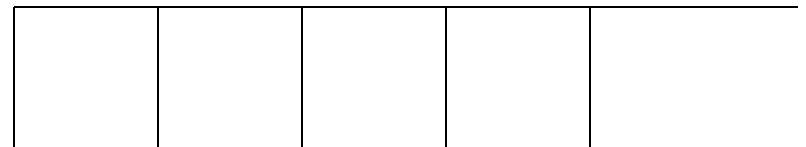
- The Nachos *Thread* object supports the following operations:
 - *Thread *Thread(char *debugName)*
 - *Fork(VoidFunctionPtr func, int arg)*
 - *void StackAllocate(VoidFunctionPtr func, int arg)*
 - *void Yield(), void Sleep(), void Finish()*

Nachos Thread's example

- Thread *th1 = new Thread("Thread1");

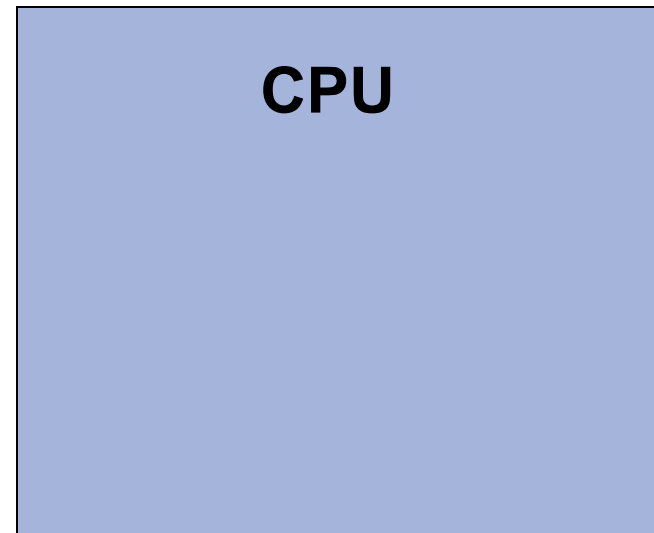
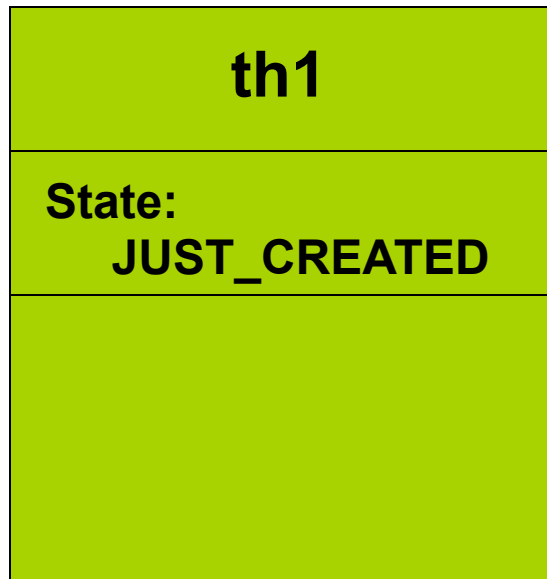


Ready list



Nachos Thread's example

- Thread *th1 = new Thread("Thread1");

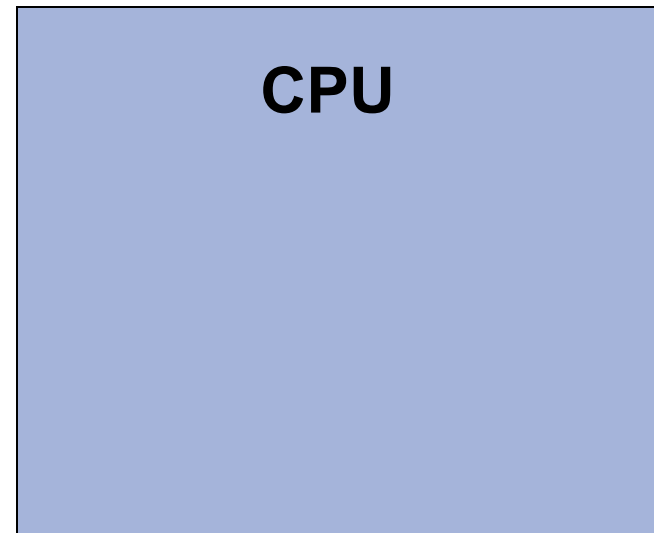
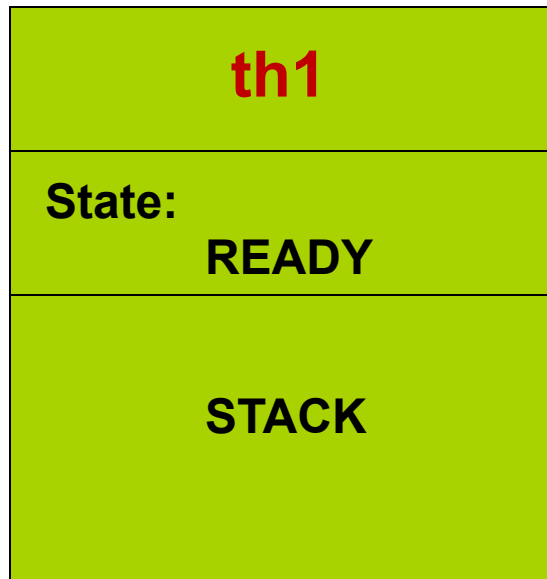


Ready list

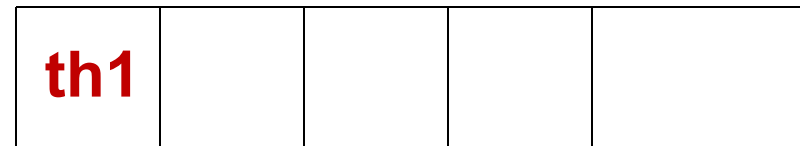


Nachos Thread's example

- th1 → Fork();

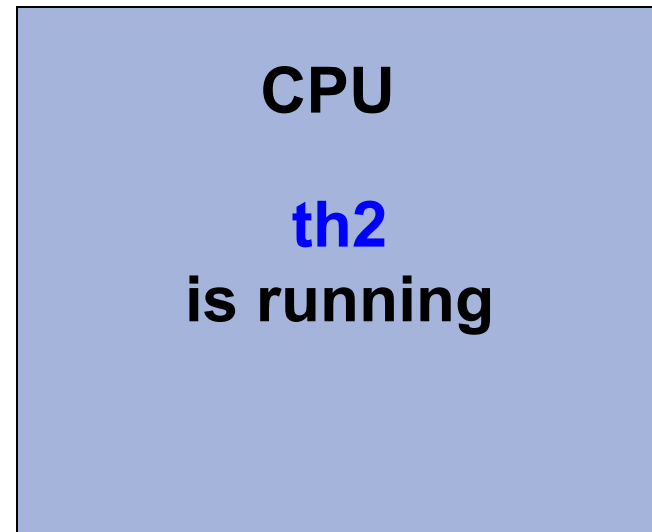
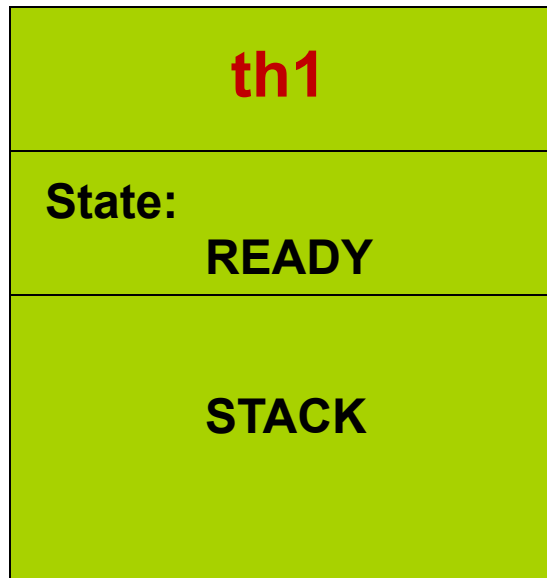


Ready list

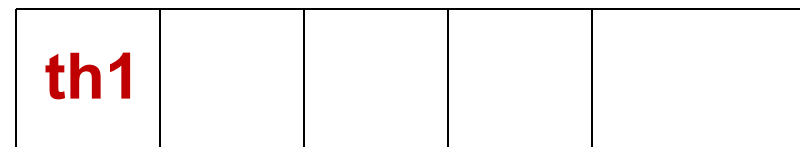


Nachos Thread's example

currentThread = **th2**;

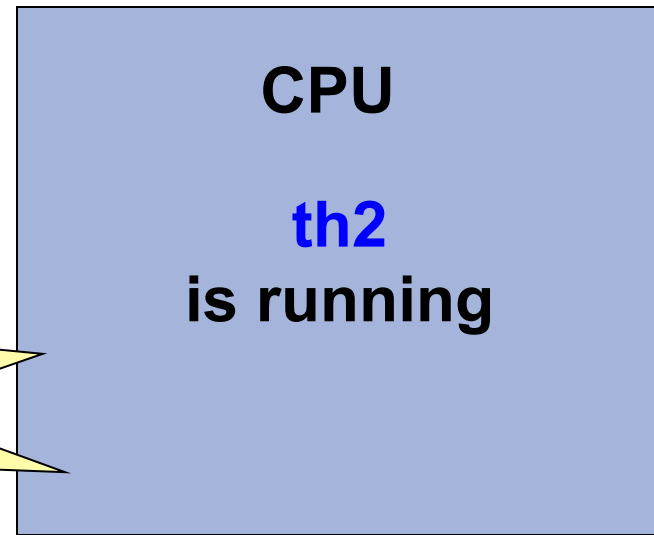
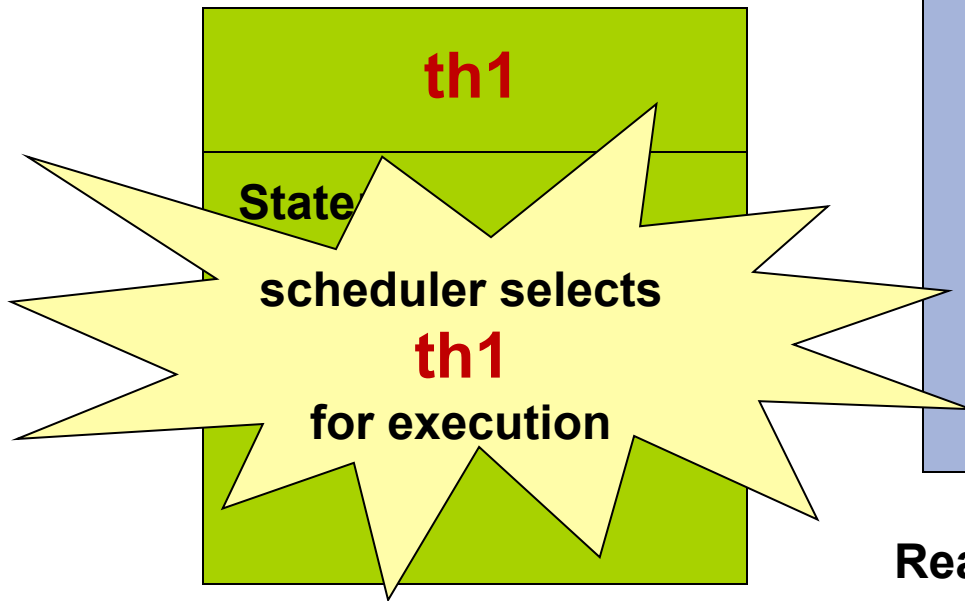


Ready list



Nachos Thread's example

currentThread = **th2**;

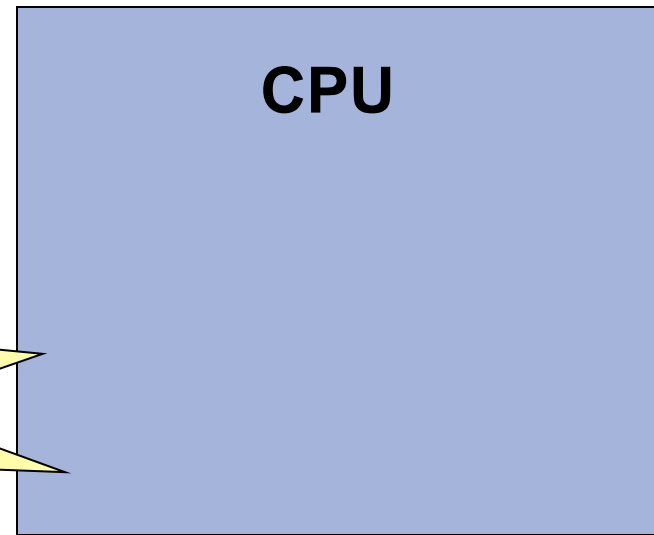
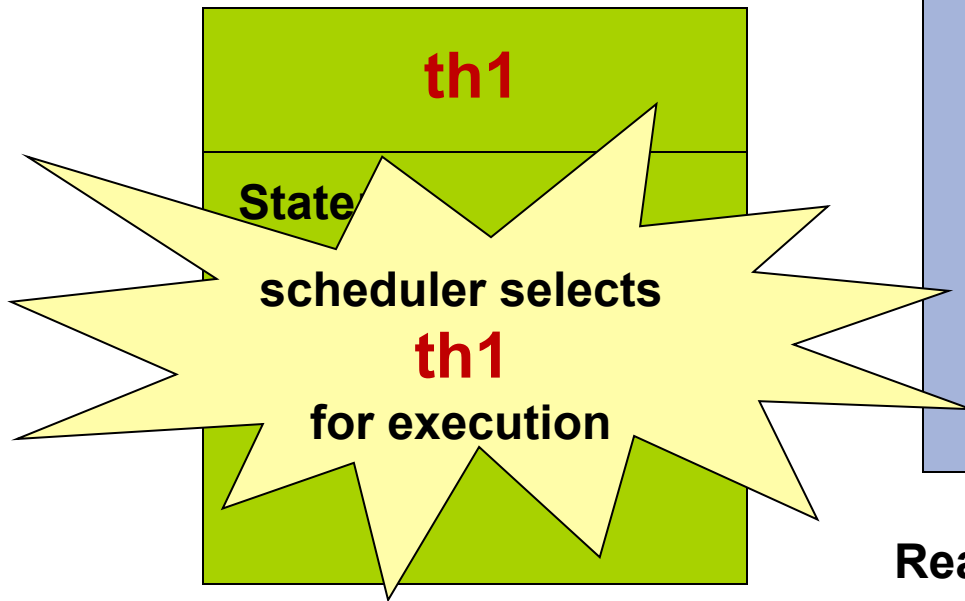


Ready list

th1				
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Nachos Thread's example

currentThread = ;

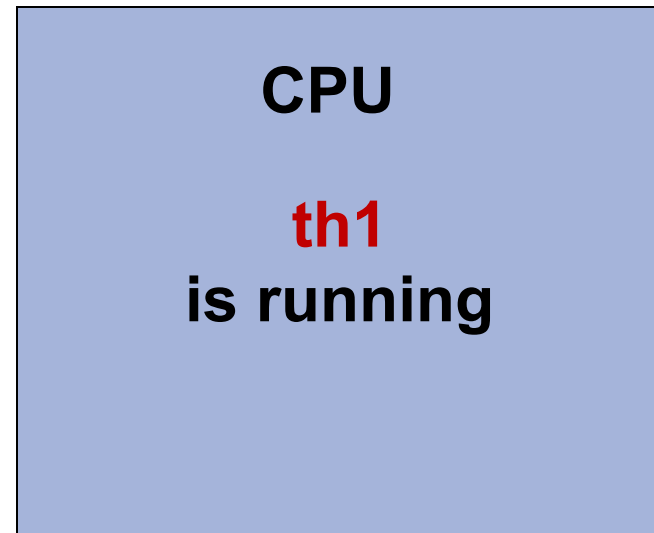
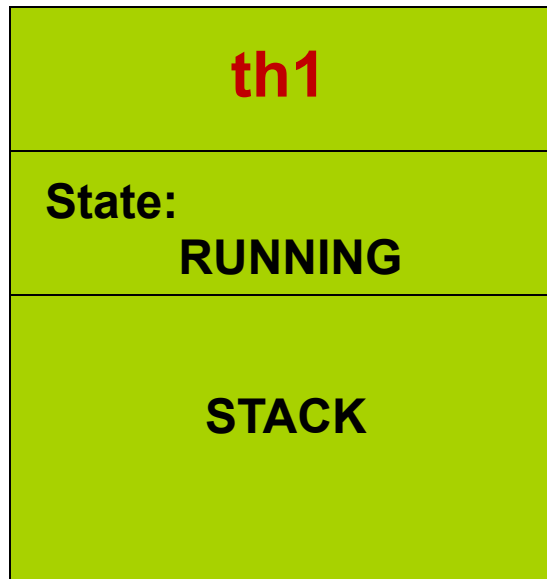


Ready list

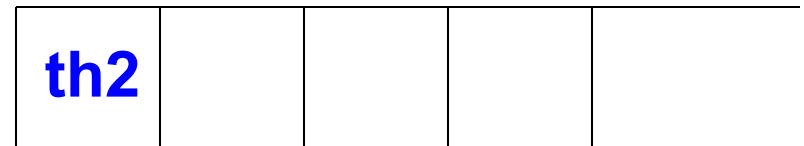
th1	th2			
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Nachos Thread's example

currentThread = **th1**;

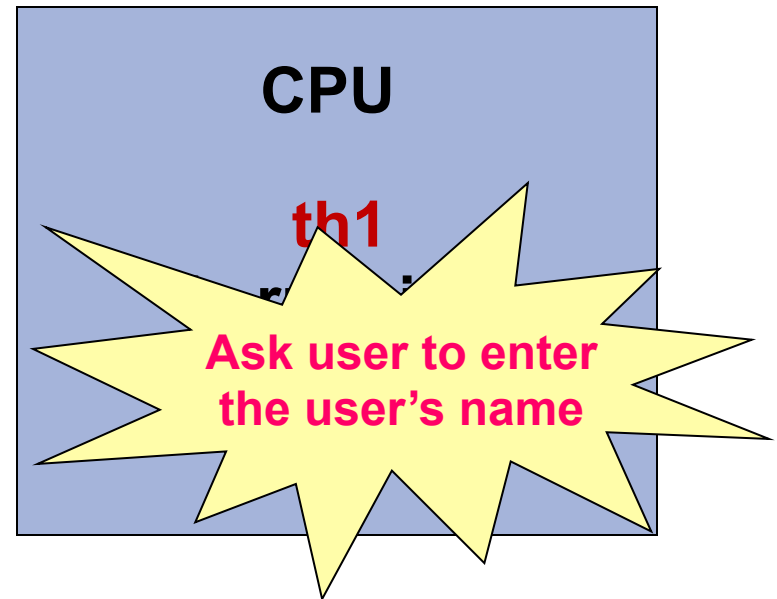
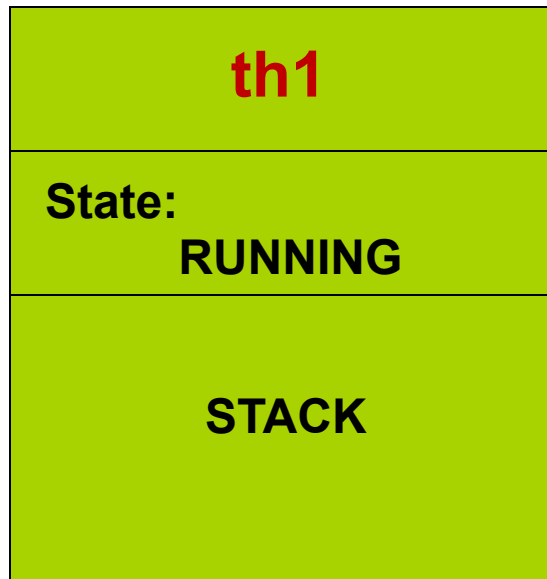


Ready list

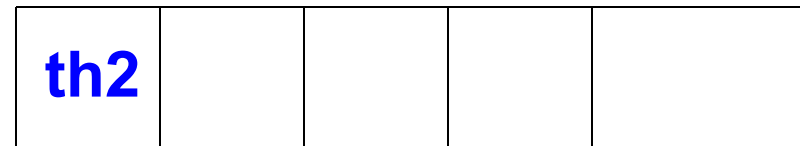


Nachos Thread's example

currentThread = **th1**;

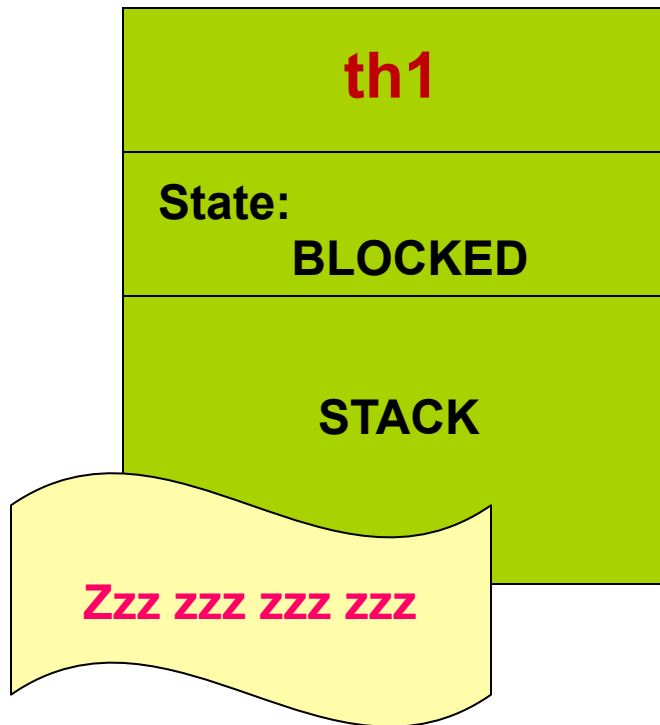


Ready list

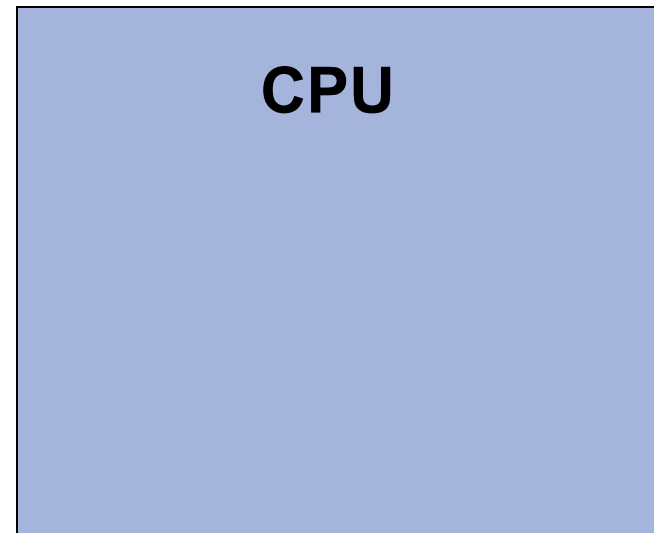


Nachos Thread's example

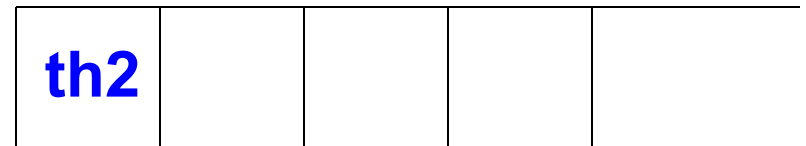
■ th1 → Sleep();



currentThread = ;



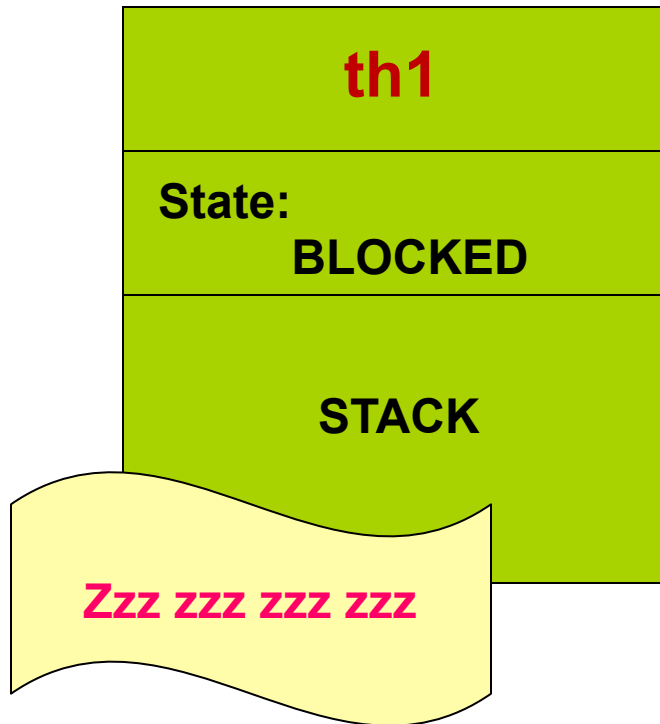
Ready list



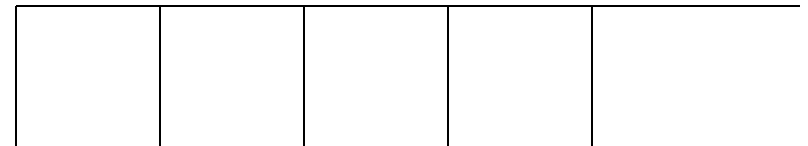
Nachos Thread's example

■ th1 → Sleep();

currentThread = **th2**;



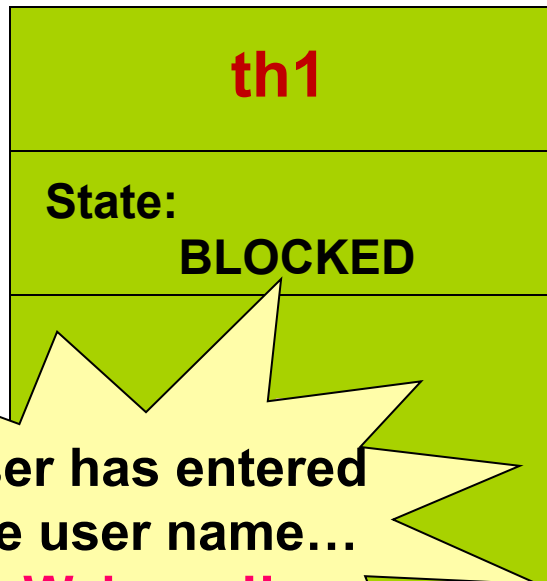
Ready list



Nachos Thread's example

■ th1 → Sleep();

currentThread = **th2**;



User has entered
the user name...
Wake up!!

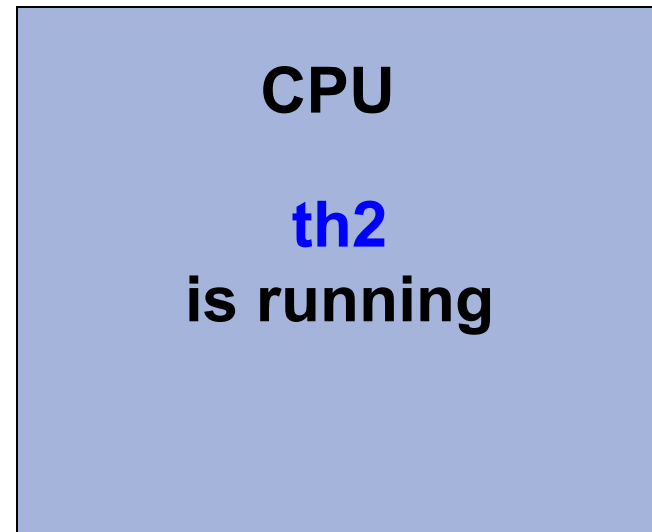
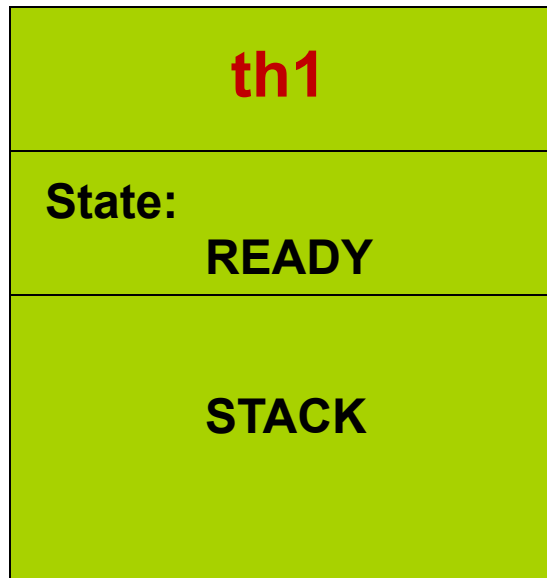


Ready list

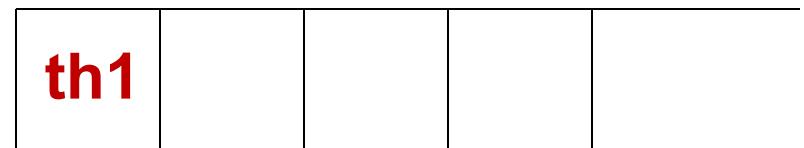


Nachos Thread's example

currentThread = **th2**;



Ready list



Nachos Thread Scheduling

- Threads that are ready to run are kept on the *ready list*.
- A process is in the **READY** state only if it has all the resources it needs, other than the CPU
- Processes **blocked** waiting for I/O, memory, etc. are generally stored in a **queue** associated with the resource being waited on.

Nachos Thread Scheduling

- The ***scheduler*** decides which thread to run next.
- The scheduler is invoked whenever the current thread wishes to give up the CPU.
 - e.g., the current thread may have initiated an I/O operation and must wait for it to complete before executing further.

Nachos Thread Scheduling

- A simple scheduling policy:
 - threads reside on a single, un-prioritized ready list, and threads are selected in a ***round-robin fashion***.
- That is, threads are always appended to the end of the ready list, and the scheduler always selects the thread **at the front of the list**.

Nachos Thread Scheduling

- Alternatively, Nachos may **preempt** the current thread in order to prevent one thread from **monopolizing** the CPU.
- Scheduling is handled by routines in the ***Scheduler*** object with the following operations:
 - *void ReadyToRun(Thread *thread)*
 - *Thread *FindNextToRun()*
 - *void Run(Thread *nextThread)*

Nachos Thread Scheduling

- `void ReadyToRun(Thread *thread)`
 - Make *thread* ready to run and place it on the ready list.
 - Note that *ReadyToRun* doesn't actually start running the thread; it simply changes its state to READY and places it on the ready list.
 - The thread won't start executing until later, when the scheduler chooses it.

Nachos Thread Scheduling

- Thread *FindNextToRun()
 - Select a ready thread and return it.
 - For round-robin fashion, *FindNextToRun* simply returns the thread at the front of the ready list.

Nachos Thread Scheduling

- `void Run(Thread *nextThread)`
 - Do the dirty work of suspending the current thread and switching to the new one.
 - Note that it is the **currently running thread** that calls *Run()*.
 - A thread calls this routine when it no longer wishes to execute.

Nachos Thread Switching

- Switching the CPU from one thread to another involves:
 - **suspending** the current thread
 - **saving** its state (e.g., registers)
 - then **restoring** the state of the thread being switched to

Nachos Thread Switching

- The thread switch actually completes at the moment a new program counter is loaded into PC.
- At that point, the CPU is no longer executing the thread switching code, it is executing code associated with the new thread.

Nachos Thread Switching

- The routine *Switch(oldThread, nextThread)* actually performs a thread switch:
 - Save all registers in *oldThread*'s *context block* and suspend it
 - load new values into the registers from the context block of the *nextThread*
 - Once the saved PC is loaded, *Switch()* is no longer executing; we are now executing instructions associated with the new thread

Nachos Thread Switching

- After returning from *Switch*, the previous thread is no longer running. Thread *nextThread* is running now.
- The routine *Switch()* is written in assembly language because it is a machine-dependent routine.
- It has to manipulate registers, look into the thread's stack, etc.