

# **COMP2611: Computer Organization**

## **MIPS Procedure**

## **MIPS procedures**

32-bit Immediate Operands

- exercises

Simple MIPS procedures

- exercises

Exercises

- ❑ Constants are frequently short and fit into 16-bit field
- ❑ But sometimes they are bigger than 16 bits, e.g. 32-bit constant
  
- ❑ **lui** (“load upper immediate”)
  - ❑ e.g. **lui reg, constant**
  - ❑ set the upper 16 bits of register **reg** to the 16-bit value specified in **constant**
  - ❑ **Set the lower 16 bits of register **reg** to zeros**
- ❑ **addi** will sign-extend the 16-immediate operand especially for negative numbers
  - ❑ For unsigned immediate
- ❑ Advisable to use **ori**
  - ❑ **For both signed and unsigned immediate**

**Question 1:** Write down the shortest sequence of MIPS instructions for the following C++ codes, assuming each variable is stored in a different register (you name it).

```
b = a + 0x37cf0010;
```

**Question 2:** Write down the shortest sequence of MIPS instructions for the following C++ codes, assuming each variable is stored in a different register (you name it).

```
b = a + 0x37cff346;
```

- ❑ Arithmetic instructions (e.g. addi, addiu): always sign extend (deem zero-extend as sign-extend for unsigned number)
- ❑ Load/store instructions (e.g. lb, lbu): always sign extend
- ❑ Logical instructions (e.g. ori, andi): always zero extend
- ❑ Set instructions (e.g. slti, sltiu): sign extend
- ❑ shift instructions (e.g. srl): always sign extend

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Exercises

## ❑ The Caller

- ❑ Puts function arguments in  $\$a0 - \$a3$  before invoking `jal`
- ❑ Pushes arguments registers ( $\$a0 - \$a3$ ), temporary registers ( $\$t0 - \$t9$ ) onto stack if needed after the call
- ❑ `jal ProcedureAddress`
  - The `jal` saves the return address which is  $(PC + 4)$  in  $\$ra$
  - Then, jump to address specified by `ProcedureAddress`
- ❑ Picks up the return values from  $\$v0 - \$v1$

## ❑ The Callee

- ❑ Pushes preserved registers ( $\$s0 - \$s8$ ), argument registers ( $\$a0 - \$a1$ ) onto stack if they are changed within callee
- ❑ Performs the procedure
- ❑ Pops the preserved registers if any from stack
- ❑ Puts up to two return results in  $\$v0 - \$v1$  if there is any
- ❑ Invokes `jr $ra` to go back to the Caller



- ❑ Since **procedures** are like small programs themselves, they may **need to use the registers**, and they **may also call other procedures** (nested calls)
  - ❑ If we don't **save** some of the values stored in the registers, they will be wiped each time we call a new procedure
- ❑ In MIPS, we need to save the registers by ourselves
- ❑ The perfect place for this is called a **stack**
  - a memory accessible only from the top (Last In First Out, LIFO)
  - placing things on the stack is called **push**
  - removing them is called **pop**
- ❑ **push** and **pop** are simply **storing** and **loading** words to and from a specific location in the memory pointed to by **the stack pointer \$sp** which always points to top of the stack

**Question 1:** Translate the following C++ function into a MIPS function, using the registers `$a0` and `$a1` for its parameters and the register `$v0` for its return value.

```
int equal(int p1, int p2) {  
    if (p1 == p2)  
        return 1;  
    return 0;  
}
```

**Question 2:** Write down the MIPS instructions that make the following call to the C++ function in the previous exercise, assuming the variable `b` is stored in the register `$s0`.

```
int b = equal(3, 4);
```

**Question 4:** The following C++ function takes as inputs the base address of an int array  $A$  and returns the minimum value in  $A$ . Using the registers  $\$a0$  and  $\$a1$  as arguments to the function,  $\$v0$  as returned value,  $\$s0$  as base address of  $A$  and  $\$s1$  as the size of  $A$ , translate the C++ function into a MIPS function

```
int minArray(int A[], int arraySize) {
    int min = A[0];
    int i = 1;
    while(i < arraySize) {
        if(min < A[i])
            min = A[i];
        i++;
    }
    return min;
}
```

## **MIPS procedures**

MIPS immediate numbers

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**Exercises**

**Question 1:** Write down the shortest sequence of MIPS instructions for the following C++ codes, assuming each variable is stored in a different register (you name it).

```
b = a + 60000;
```

**Solution:**

```
addi $s0, $s1, 60000 # $s0 and $s1 store b and a  
                        respectively
```

**Question 2:**

```
void saveElement(int a[], int x) {  
    a[x] = x;  
}
```

Translate the above C++ function into a MIPS function, assuming the registers `$a0` and `$a1` store the parameters. `$s0` is the only extra register that can be used inside your function. The stack can also be used. Your function must work for the following MIPS sequence of calls to it.

```
la $a0, list1      #assuming an array list1 is already defined  
addi $a1, $s0, 0  
jal saveElement  
addi $a1, $s0, 1  
jal saveElement
```

**Solution to Question 2:**

```
saveElement: addi $sp, $sp, -4
              sw  $a1, 0($sp)
              sll $s0, $a1, 2
              add $s0, $s0, $a0
              sw  $a1, 0($s0)
              lw  $s0, 0($sp)
              addi $sp, $sp, 4
              jr  $ra
```